COMPLEXITY AND ALGORITHMS FOR COMPUTING VORONOI CELLS OF LATTICES

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ABSTRACT. In this paper we are concerned with finding the vertices of the Voronoi cell of a Euclidean lattice. Given a basis of a lattice, we prove that computing the number of vertices is a $\#$P-hard problem. On the other hand we describe an algorithm for this problem which is especially suited for low dimensional (say dimensions at most 12) and for highly-symmetric lattices. We use our implementation, which drastically outperforms those of current computer algebra systems, to find the vertices of Voronoi cells and quantizer constants of some prominent lattices.

1. INTRODUCTION

Let $L = B\mathbb{Z}^m \subseteq \mathbb{R}^n$ be a *lattice* of rank m in Euclidean space given by a matrix $B \in \mathbb{R}^{n \times m}$ of rank m. By $\lim L$ we denote the linear subspace spanned by the elements of L. The *Voronoi cell* of L is

$$\mathrm{V}(L) = \{x \in \lim L : \|x\| \leq \|x - v\| \text{ for all } v \in L\}.$$

The Voronoi cell of a lattice is a centrally symmetric, convex polytope. The polytopes $V(L) + v$ for $v \in L$ tile $\lim L$. The study of Voronoi cells is motivated by the fact that most important geometric lattice parameters have a direct interpretation in terms of the Voronoi cell: The *determinant* $\det L$ equals the volume of $\mathrm{V}(L)$, the *packing radius* $\lambda(L)$ equals the inradius of $\mathrm{V}(L)$, the *covering radius* $\mu(L)$ equals the circumradius of $\mathrm{V}(L)$, and the *quantizer constant* $G(L)$ is

$$G(L) = (\det L)^{-(1+2/n)} \int_{\mathrm{V}(L)} \|x\|^2 dx.$$

In this paper we consider theoretical and practical aspects of the computation of the covering radius as well as the quantizer constant of a lattice. These two parameters have many applications, we just name a few: By computing the covering radius, we get an upper bound for the lattice sphere covering problem, which is the problem of minimizing the covering radius among the lattices of fixed determinant

(see [CS99, Chapter 2] and [SV06]). The computation of the covering radius of the Leech lattice in [CS99, Chapter 23] had a major impact on the study of hyperbolic reflection groups (see [CS99, Chapter 27]). An upper bound for the Frobenius number of a set of integers can be obtained from the covering radius of a suitable lattice (see [FR05]). A recent application comes from public key cryptography; Micciancio [Mic04] found a new connection between the average-case complexity of finding the packing radius and the worst-case complexity of determining the covering radius. In information theory, the quality of a lattice as a vector quantizer is determined by its quantizer constant (see [GG92]).

The structure of this paper is as follows. In Section 2 we discuss the computational complexity of the covering radius problem. We prove that the related problem of counting vertices of the Voronoi cell is $\#$P-hard. As a byproduct of our construction, we show that the lattice isomorphism problem is at least as difficult as the graph isomorphism problem. We turn to practical algorithms for the covering radius problem in Section 3. There we describe an algorithm which computes the vertices of the Voronoi cell of a lattice. Based on this algorithm we give an algorithm for computing the quantizer constant in Section 4. In Section 5 we report on computations with our implementation. We determine the exact covering radius and quantizer constants of many prominent lattices which were not known before.

2. COMPUTATIONAL COMPLEXITY

We formulate the covering radius problem as a decision problem.

Problem 1. *Covering radius problem*

> **Input:** m, n, $B \in \mathbb{Q}^{m \times n}$, $\mu \in \mathbb{Q}$.
> **Output:** Yes, *if* $\mu(B\mathbb{Z}^n) \leq \mu$, No *otherwise.*

It is conjectured (see [Mic04, Section 1.1]) that the covering radius problem is NP-hard. Haviv and Regev [HR06] showed that there is a constant c_p so that the covering radius in the l_p-norm is Π_2-hard to approximate within a constant less than c_p for any large enough p. In [GMR05] Guruswami, Micciancio and Regev proved that approximating it within a factor of $O(\sqrt{m/\log m})$ for a lattice of rank m cannot be NP-hard unless the polynomial hierarchy collapses.

Currently, there is only one known general and practical method to compute $\mu(L)$ for a lattice L: First one enumerates the vertices of $\mathrm{V}(L)$ and then one finds the vertex with largest norm. The number of vertices of $\mathrm{V}(L)$ can be as large as $(m + 1)!$ and furthermore, as we show in Theorem 1, even computing this number is $\#$P-hard.

Problem 2. *Vertices of a lattice Voronoi cell*

> **Input:** m, n, $B \in \mathbb{Q}^{m \times n}$.
> **Output:** *Number of vertices of* $\mathrm{V}(B\mathbb{Z}^n)$.

Theorem 1. *The problem "Vertices of a lattice Voronoi cell" is $\#$P-hard.*

It will be obvious from the proof that we could restrict the problem to the case $m = n$. We reduce the problem "Acyclic orientations of a graph", which Linial [Lin86] showed to be $\#\mathrm{P}$-complete, to Problem 2.

Problem 3. *Acyclic orientations of a graph*
   Input: *A graph $G = (V, E)$.*
   Output: *The number of orientations of G with no directed circuit.*

The structure of the proof of Theorem 1 is as follows: In Section 2.1 we construct a matrix B with columns indexed by E defining a lattice $L(G) = B\mathbb{Z}^E$ from G in polynomial time. Then we show that the vertices of the Voronoi cell of $V(L(G))$ are in bijection with the acyclic orientations of G. To establish this bijection we need several intermediate steps. In Section 2.2 we associate to G a hyperplane arrangement $H(G)$ whose full-dimensional regions are in bijection with the acyclic orientations of G. In Section 2.3 we recall that the full-dimensional regions of a hyperplane arrangement are in bijection with the vertices of a zonotope associated to the hyperplane arrangement. These two steps are standard and we cover them rather briefly. In Section 2.4 we show that the Voronoi cell of $L(G)$ is a zonotope which, up to a linear transformation, is the one associated to the hyperplane arrangement $H(G)$. In Section 2.5, as a byproduct of this construction, we show that the lattice isomorphism problem is at least as difficult as the graph isomorphism problem.

2.1. **From graphs to lattices.** Let $G = (V, E)$ be a connected graph with vertex set $V = \{1, \ldots, n\}$ and edge set E. We consider the following orientation of the edges of G: The head of an edge $e = \{v, w\} \in E$ is $e^+ = \max\{v, w\}$ and the tail is $e^- = \min\{v, w\}$.

Let $T \subseteq E$ be the edge set of a spanning tree of G, and let $e \in T$. Deleting e from T divides T into two connected components with vertex sets T_e^+ and T_e^-, where $e^+ \in T_e^+$ and $e^- \in T_e^-$. Define the vector $b_{T,e} \in \mathbb{Z}^E$ by

$$b_{T,e}(f) = \begin{cases} 1, & \text{if } f^+ \in T_e^+ \text{ and } f^- \in T_e^-, \\ -1, & \text{if } f^- \in T_e^+ \text{ and } f^+ \in T_e^-, \\ 0, & \text{otherwise.} \end{cases}$$

Then

$$L(G, T) = \left\{ \sum_{e \in T} \alpha_e b_{T,e} : \alpha_e \in \mathbb{Z} \right\} \subseteq \mathbb{Z}^E$$

is a lattice of rank $n - 1$.

Proposition 1. *Let T and T' be spanning trees of G. Then, $L(G, T) = L(G, T')$.*

Proof. Since one can connect any two spanning trees by a sequence of transformations of the form $T \leftrightarrow T \setminus \{e\} \cup \{f\}$ it suffices to prove the proposition for $T' = T \setminus \{e\} \cup \{f\}$. Let $g \in T'$. If $g = f$, then $b_{T',f} = \pm b_{T,e}$. If $g \in T$, then it is easy to see that $b_{T',g} = b_{T,g} + \alpha b_{T,e}$ with $\alpha \in \{-1, 0, +1\}$. $\square$

In the following we omit the spanning tree T from the notation $L(G, T)$ and just write $L(G)$. Note that one can find a basis of $L(G)$ given G in polynomial time.

2.2. **From graphs to hyperplane arrangements.** A matrix

$$(1) \qquad\qquad V = (v_1, \ldots, v_m) \in \mathbb{R}^{n \times m}$$

with non-zero column vectors $v_i \in \mathbb{R}^n$ gives an *arrangement of hyperplanes*

$$H(V) = \{H_1, \ldots, H_m\} \quad \text{with} \quad H_i = \{c \in \mathbb{R}^n : c \cdot v_i = 0\}.$$

The hyperplane arrangement $H(V)$ divides the space $\mathbb{R}^n$ into polyhedral cones, called regions, of different dimensions. The regions are partially ordered by inclusion.

To associate a hyperplane arrangement $H(G)$ with G we consider the incidence matrix $D_G \in \mathbb{R}^{V \times E}$ of G which is given by

$$D_G(v, e) = \begin{cases} 1, & \text{if } v = e^+, \\ -1, & \text{if } v = e^-, \\ 0, & \text{otherwise.} \end{cases}$$

Then we define the hyperplane arrangement of G by $H(G) = H(D_G)$.

In [GZ83, Lemma 7.1] Greene and Zaslavsky show that the full-dimensional regions of $H(G)$ are in bijection with the acyclic orientations of G: Let $\vec{E}$ be an acyclic orientation of E. Then a full-dimensional region of $H(G)$ is given by

$$\mathrm{Reg}(\vec{E}) = \{x \in \mathbb{R}^V : x_v < x_w \text{ if } (v, w) \in \vec{E}\}.$$

Let R be a full-dimensional region of $H(G)$. Then an acyclic orientation of E is given by

$$\vec{E}(R) = \{(v, w) : \{v, w\} \in E \text{ and } x_v < x_w \text{ if } x \in R\}.$$

Obviously, $\mathrm{Reg}(\vec{E}(R)) = R$.

2.3. **Hyperplane arrangements and zonotopes.** The matrix V in (1) defines a *zonotope* $Z(V)$ by

$$Z(V) = \left\{ \sum_{i=1}^m \alpha_i v_i : -1 \le \alpha_i \le 1 \right\}.$$

The faces of $Z(V)$ are partially ordered by inclusion. It is a well-known fact (see e.g. [Zie95, Theorem 7.16]) that the partially ordered set of regions of the hyperplane arrangement $H(V)$ is anti-isomorphic to the partially ordered set of faces of $Z(V)$: Let R be a region of $H(V)$. Then the corresponding face of $Z(V)$ is given by

$$\mathrm{Face}(R) = \{y \in Z(V) : x \cdot y = \max_{z \in Z(V)} x \cdot z\},$$

where $\mathrm{Face}(R)$ does not depend on the choice of x lying in R. Let F be a face of $Z(V)$. Then the corresponding region of $H(V)$ is given by

$$\mathrm{Reg}(F) = \{x \in \mathbb{R}^n : \max_{z \in Z(V)} x \cdot z = x \cdot y\},$$

where $\mathrm{Reg}(F)$ does not depend on the choice of y in the relative interior of F. Obviously, $\mathrm{Face}(\mathrm{Reg}(F)) = F$ and $F' \subseteq F$ if and only if $\mathrm{Reg}(F') \supseteq \mathrm{Reg}(F)$. In

particular, the full-dimensional regions of $H(V)$ are in bijection with the vertices of $Z(V)$.

2.4. **From lattices to zonotopes.** Let $L \subseteq \mathbb{R}^n$ be a lattice. The *support* of a vector $v \in L$ is $\underline{v} = \{i \in \{1, \ldots, n\} : v_i \neq 0\}$. It is called *elementary* if $v \in \{-1, 0, +1\}^n \setminus \{0\}$ and if v has minimal support among all vectors in $L \setminus \{0\}$. We say that two vectors $v, w \in L$ are *conformal* if $v_i w_i \geq 0$ for all $i = 1, \ldots, n$. The lattice L is called *regular* if for every vector $v \in L \setminus \{0\}$ there exists an elementary vector $u \in L$ with $\underline{u} \subseteq \underline{v}$.

Lemma 1. ([Tut71, Chapter 1])

(i) For any graph G the lattice $L(G)$ is regular.

(ii) If L is a regular lattice, then every $v \in L$ can be written as a sum of pairwise conformal elementary vectors.

(iii) If L is a regular lattice, $v \in L$ is elementary, and $u \in L$ satisfies $\underline{u} = \underline{v}$, then there exists a factor $\alpha \in \mathbb{Z}$ such that $u = \alpha v$.

A vector $v \in L$ for which $\mathrm{V}(L) \cap \{x \in \mathbb{R}^n : x \cdot v = \frac{1}{2} v \cdot v\}$ is a facet of $\mathrm{V}(L)$ is called *relevant*. Voronoi characterizes in [Vor08, page 277] the relevant vectors of L: A nonzero vector $v \in L$ is relevant if and only if $\pm v$ are the only shortest vectors in $v + 2L$.

Proposition 2. *In a regular lattice, a vector is elementary if and only if it is relevant.*

Proof. Let $v \in L$ be a relevant vector. By Lemma 1 (ii), we can write $v = \sum_{k=1}^m w_k$ as a sum of pairwise conformal elementary vectors $w_k \in L$. Assume that $m \geq 2$. Defining $u = v - 2w_1$ gives $u \neq \pm v$ and $u \cdot u = v \cdot v - 4(v - w_1) \cdot w_1$. Since the vectors w_k, $k = 1, \ldots, m$, are pairwise conformal we have $(v - w_1) \cdot w_1 \geq 0$, and $\pm v$ is not the unique shortest vector in $v + 2L$. In this case v cannot be a relevant vector. Hence, $m = 1$ and v is an elementary vector.

Let $v \in L$ be an elementary vector, and let $u \in v + 2L$ be a lattice vector with $u \neq \pm v$. We have $v - u \in 2L \subseteq 2\mathbb{Z}^n$ and $v_i \in \{-1, 0, +1\}$, which shows $\underline{v} \subseteq \underline{u}$. The case $\underline{v} \neq \underline{u}$ immediately leads to $v \cdot v < u \cdot u$. If $\underline{v} = \underline{u}$, then by Lemma 1 (iii), there exists a factor $\alpha \in \mathbb{Z} \setminus \{-1, +1\}$ so that $u = \alpha v$, hence $v \cdot v < u \cdot u$. In both cases $\pm v$ are the only shortest vectors in $v + 2L$. Hence, v is a relevant vector. $\square$

The following special case of Farkas lemma is proved e.g. in [Roc70, Theorem 22.6].

Lemma 2. *Let $L \subseteq \mathbb{R}^n$ be a regular lattice. Let $x \in \mathbb{R}^n$ be a vector, and let $\alpha_1, \ldots, \alpha_n \in \mathbb{R} \cup \{\pm\infty\}$. Either there exists a vector $y' \in (\operatorname{lin} L)^\perp$ lying in $x + \prod_{i=1}^n [-\alpha_i, \alpha_i]$, or there exists a vector $y \in \operatorname{lin} L$ such that for all $z \in x + \prod_{i=1}^n [-\alpha_i, \alpha_i]$ the inequality $y \cdot z > 0$ holds. If the second condition holds, then one can choose y to be an elementary vector of L.*

Theorem 2. *Let $L \subseteq \mathbb{R}^n$ be a regular lattice. Let $P \in \mathbb{R}^{n \times n}$ be the matrix of the orthogonal projection of $\mathbb{R}^n$ onto $\operatorname{lin} L$. Then, $\mathrm{V}(L) = \frac{1}{2} Z(P) = P([-1/2, 1/2]^n)$.*

Proof. Suppose that $x \in [-1/2, 1/2]^n$. For all $v \in \mathbb{Z}^n \setminus \{0\}$ the inequality $x \cdot v \leq \frac{1}{2} v \cdot v$ holds. Write $x = y + y'$ with $y = Px \in \lin L$ and $y' \in (\lin L)^\perp$. For all $v \in L \setminus \{0\}$ we have $y \cdot v = x \cdot v - y' \cdot v \leq \frac{1}{2} v \cdot v$. Thus, $Px \in \mathrm{V}(L)$.

Suppose now that $y \in \mathrm{V}(L)$. If there exists $x \in (-y + [-1/2, 1/2]^n) \cap (\lin L)^\perp$, then $y + x \in [-1/2, 1/2]^n$ and $P(y + x) = y$. Assume that such a vector does not exist. Then by Lemma 2 there is an elementary lattice vector $v \in L$ so that $v \cdot (-y + [-1/2, 1/2]^n) > 0$. This implies $v \cdot (-y - \frac{1}{2}v) > 0$. Hence, $-y \notin \mathrm{V}(L)$. Since $\mathrm{V}(L)$ is centrally symmetric, this contradicts the assumption $y \in \mathrm{V}(L)$. □

In [Big97, Proposition 8.1] Biggs shows that for the lattice $L(G)$ the matrix P can be written in the form $P = XD_G$ where $D_G \in \mathbb{R}^{V \times E}$ is the incidence matrix of G and $X \in \mathbb{R}^{E \times V}$ is given by

$$(2) \qquad X(e, v) = \frac{\text{number of spanning trees } T \text{ with } e \in T \text{ and } v \in T_e^+}{\text{number of spanning trees of } G}.$$

Furthermore, the linear map given by X restricted to the image of D_G is a bijection.

Thus, the zonotope $Z(P)$ which is the Voronoi cell of $L(G)$ equals $\frac{1}{2}XZ(D_G)$. Hence, there is a linear isomorphism between the faces of $\mathrm{V}(L(G))$ and those of $Z(D_G)$. This completes the proof of Theorem 1.

Using a straightforward computation we get the following proposition.

Proposition 3. *Using the notation in* (2)*, the covering radius of the lattice $L(G)$ is given by*

$$(3) \quad \mu(L(G))^2 = \max_{x \in [-1/2, 1/2]^E} \sum_{e \in E} \left(\sum_{f \in E} (X(e, f^+) - X(e, f^-))x(f) \right)^2.$$

Unfortunately, we do not have a combinatorial interpretation of (3). Finding one could lead to a proof of the NP-hardness of the covering radius problem.

2.5. **Lattice isomorphism problem.** Using the construction $L(G)$ used in the proof of Theorem 1, we reduce the graph isomorphism problem to the lattice isomorphism problem in polynomial time. We don't know whether one can give a reverse polynomial time reduction. For the graph isomorphism problem no polynomial time algorithm is known. It is generally believed to lie in NP ∩ co-NP. So it is unlikely that it is NP-hard. For more information on the computational complexity of this problem, see the book [KST93] of Köbler, Schöning and Tóran.

Problem 4. *Lattice isomorphism problem*

> **Input:** $m, n, B, B' \in \mathbb{Q}^{m \times n}$ *matrices of rank m.*
> **Output:** Yes*, if there is an orthogonal transformation O so that $OB\mathbb{Z}^n = B'\mathbb{Z}^n$,* No *otherwise.*

Problem 5. *Graph isomorphism problem*

> **Input:** *Graphs $G = (V, E_G)$, $H = (V, E_H)$.*
> **Output:** Yes*, if there is a permutation $\sigma : V \to V$ so that for all $v, w \in V$ we have $\{v, w\} \in E_G$ if and only if $\{\sigma(v), \sigma(w)\} \in E_H$,* No *otherwise.*

Theorem 3. *There is a polynomial time reduction of the graph isomorphism problem to the lattice isomorphism problem.*

Proof. Let $G = (V, E_G)$ and $H = (V, E_H)$ be graphs. We modify G and H by adding three extra vertices to V each adjacent to all vertices in V. We call the new graphs G' and H' which are by construction 3-connected and they are isomorphic if and only if G and H are isomorphic.

It is clear that the lattices $L(G')$ and $L(H')$ are isomorphic whenever G' and H' are. For this direction it would be enough to consider the original graphs G and H.

Now suppose that the lattices $L(G')$ and $L(H')$ are isomorphic. We apply the 2-Isomorphism-Theorem of Whitney (actually we only use the easy subcase of 3-connected graphs [Oxl92, Lemma 5.3.2]): Because the graphs G' and H' are 3-connected and there is a bijection between the elementary vectors preserving conformality, the graphs G' and H' are isomorphic. $\qquad\square$

3. ALGORITHMS

In this section we describe an algorithm which computes all vertices of a lattice Voronoi cell. Our focus is on implementability and practical performance, using the symmetries of the lattice. In fact, the algorithm computes all full-dimensional Delone cells and the adjacencies between them up to equivalence. We give necessary definitions in Section 3.1. In Section 3.2 we describe the algorithm's main steps and in the following sections we give details about its subalgorithms. In Section 3.6, we explain how to use Gram matrices instead of lattice basis and in Section 3.7 we compare our method with existing algorithms.

3.1. **Notation.** From now on, we assume lattices $L \subseteq \mathbb{R}^n$ to have full rank n.

To encode the vertices of $\mathrm{V}(L)$ we use Delone cells. A point $x \in \mathbb{R}^n$ defines a *Delone cell* $D(x)$ by

$$D(x) = \mathrm{conv}\{v \in L : \|x - v\| = \min_{w \in L} \|x - w\|\}.$$

Denote by $S(x, r)$ the sphere with center x and radius r. For $r = \min_{v \in L} \|x - v\|$, the sphere $S(x, r)$ is called *empty*, since there is no lattice point inside. In this case the polytope $D(x)$ is the convex hull of $S(x, r) \cap L$. The Delone cell of a vertex of $\mathrm{V}(L)$ is characterized among all Delone cells by the following properties: The origin is a vertex of $D(x)$ and $D(x)$ is full-dimensional.

It is well known (see e.g. [Ede01]) that the Delone cells are the projections of the faces of the infinite $(n + 1)$-dimensional polyhedral set

$$\mathrm{Lift}(L) = \mathrm{conv}\{(x, \|x\|^2) : x \in L\}.$$

The task of finding a vertex of a Delone cell of a point x, given a lattice basis of L, is called the closest vector problem. Generally this is an NP-hard problem [DKS03]; however, there are algorithms and implementations available which can solve this problem rather efficiently in low dimensions.

The *orthogonal group* $\mathrm{O}(L)$ of L is the group of all orthogonal transformations $A \in \mathrm{O}(\mathbb{R}^n)$ fixing L, i.e. $A(L) = L$. The *isometry group* $\mathrm{Iso}(L)$ of L is the group

generated by $\mathrm{O}(L)$ and all lattice translations $t_v : \mathbb{R}^n \to \mathbb{R}^n$ with $t_v(x) = x + v$ for $v \in L$.

We say that two vertices x and x' of $\mathrm{V}(L)$ are equivalent if there is an $A \in \mathrm{O}(L)$ so that $A(x) = x'$. Correspondingly, we say that two Delone cells $D(x)$ and $D(x')$ are equivalent if there is an $A \in \mathrm{Iso}(L)$ so that $A(D(x)) = D(x')$.

3.2. **Main algorithm.** Our algorithm finds a complete list of inequivalent full-dimensional Delone cells of L with respect to $\mathrm{Iso}(L)$. The enumeration process is a graph traversal algorithm of the graph of equivalence classes of full-dimensional Delone cells of L. Two equivalence classes are adjacent whenever there is a facet between two of its representatives. Note that this graph can have loops and multiple edges.

For the graph traversal algorithm below one needs four subalgorithms, which we explain in the following sections.

Input: n, $B \in \mathbb{Q}^{n \times n}$ matrix of rank n.
Output: Set $\mathcal{M}$ of all inequivalent full-dimensional Delone cells of the lattice $B\mathbb{Z}^n$ with respect to the group $\mathrm{Iso}(B\mathbb{Z}^n)$.

$x \leftarrow$ an initial vertex of $\mathrm{V}(B\mathbb{Z}^n)$.   (Section 3.3)
$T \leftarrow \{D(x)\}$.
$\mathcal{M} \leftarrow \emptyset$.
while there is a $D \in T$ **do**
    $\mathcal{M} \leftarrow \mathcal{M} \cup \{D\}$.
    $T \leftarrow T \setminus \{D\}$.
    $\mathcal{F} \leftarrow$ facets of D.   (Section 3.4)
    for $F \in \mathcal{F}$ **do**
        $D' \leftarrow$ full-dimensional Delone cell with $F = D \cap D'$.   (Section 3.4)
        if D' is not equivalent to a Delone cell in $\mathcal{M} \cup T$ **then**   (Section 3.5)
            $T \leftarrow T \cup \{D'\}$.
        end if
    end for
end while

Two full-dimensional Delone cells $D(x)$ and $v + D(x)$, both containing the origin, are equivalent under $\mathrm{O}(L)$ if and only if 0 and $-v$ are equivalent under the stabilizer group of $D(x)$ in $\mathrm{Iso}(L)$. As a consequence, we can compute the vertices of $\mathrm{V}(L)$ under $\mathrm{O}(L)$ in the following way: For every orbit of full-dimensional Delone cells given by a representative $D(x)$, we compute the orbits of vertices of $D(x)$ under the stabilizer group and get the corresponding orbits of vertices of $\mathrm{V}(L)$ under $\mathrm{O}(L)$.

3.3. **Finding an initial vertex.** Now we explain a method for computing an initial vertex of the Voronoi cell of a lattice, i.e. a full-dimensional Delone cell containing the origin. The method we propose is a so-called cutting-plane algorithm, which is a well-known technique in combinatorial optimization.

Let us describe the geometric idea. We start with an outer approximation of the Voronoi cell given by linear inequalities. The first outer approximation is the
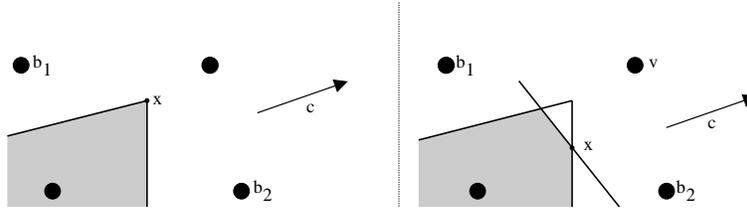
FIGURE 1. Finding an initial vertex of $V(L)$

polytope defined by the inequalities $\pm b_i \cdot x \leq \frac{1}{2} b_i \cdot b_i$ for given lattice basis vectors $b_1, \ldots, b_n$. Then we find a vertex x of the approximation by linear programming (see e.g. [Sch86]). Deciding whether the vertex x belongs to the Voronoi cell $V(L)$ can be done as follows: Compute the vertices of the Delone cell $D(x)$. If the origin is a vertex of $D(x)$, then x is a vertex of $V(L)$. Otherwise x is not contained in $V(L)$, and for all vertices v of $D(x)$ we have the strict inequality $\|x - v\| < \|v\|$. So the new linear inequalities $v \cdot x \leq \frac{1}{2} v \cdot v$ together with the old ones provide a tighter outer approximation of the Voronoi cell. Since we started with a compact outer approximation, finitely many iterations of these steps suffice to find a vertex of the Voronoi cell.

One advantage of this method is that the computation of all facets of the Voronoi cell is not required, i.e. we do not use Voronoi's characterization [Vor08, page 277] of facet defining vectors, which involves solving exponentially many closest vector problems. Figure 1 illustrates this algorithm.

Input: n, $B = (b_1, \ldots, b_n) \in \mathbb{Q}^{n \times n}$ matrix of rank n.
Output: vertex x of $V(B\mathbb{Z}^n)$.

$c \leftarrow$ random vector in $\mathbb{Q}^n$.
$\mathcal{B} \leftarrow \{\pm b_1, \ldots, \pm b_n\}$.
do
    $x \leftarrow$ a vertex of the polytope $\{x : b \cdot x \leq \frac{1}{2} b \cdot b$ for all $b \in \mathcal{B}\}$,
        which maximizes $c \cdot x$.
    $\mathcal{E} \leftarrow$ set of closest lattice vectors in $B\mathbb{Z}^n$ to x.
    if $0 \in \mathcal{E}$ **then**
        return x.
    end if
    $\mathcal{B} \leftarrow \mathcal{B} \cup \mathcal{E}$.
end do

3.4. **Computing facets of, and finding adjacent Delone cells.** We want to determine the facets of a full-dimensional Delone cell, which is given by its vertex set. This representation conversion problem can be solved by many different methods. For details and implementations we refer to `cdd` [Fuk95], `lrs` [Avi93], `pd` [Mar97] and `porta` [CL97].

In order to exploit the symmetries we use the *adjacency decomposition method* (see [CR96, BDS07, DSV07]). It allows to compute a complete list of inequivalent
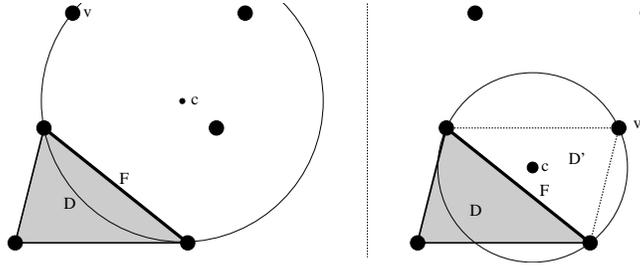
FIGURE 2. Finding D', the full-dimensional Delone cell adjacent to D at F

facet representatives: We compute an initial facet by linear programming and insert it into the list of orbit representatives of facets. From any such orbit, we compute the list of facets adjacent to a representative and insert it, if necessary, into the list of representatives until all orbits have been treated. Computing adjacent facets is itself a representation conversion problem in one dimension lower. So this method can be applied recursively (see [BDS07, DSV07]). Note that our main algorithm is itself an adjacency decomposition method.

After the computation of facets, we can compute adjacent full-dimensional Delone cells: We take an initial vertex v and so get a tentative empty sphere. If the sphere is not empty, then we find another vertex v and iterate until the sphere is indeed empty. Figure 2 illustrates this algorithm.

Input: n, $B \in \mathbb{Q}^{n \times n}$ matrix of rank n, a full-dimensional Delone cell D and a facet F of D.
Output: vertex set $\mathcal{V}'$ of a full-dimensional Delone cell D' with $D \cap D' = F$.

$\phi \leftarrow$ affine function on $\mathbb{R}^n$ with $F = \{x \in D \ : \ \phi(x) = 0\}$ and $\phi(x) > 0$
    on $D - F$.
$\mathcal{V}_F \leftarrow$ vertices of D belonging to F.
$v \leftarrow$ a point of $B\mathbb{Z}^n$ with $\phi(v) < 0$.
do
    $S(c, r) \leftarrow$ sphere around $\mathcal{V}_F \cup \{v\}$.
    $\mathcal{V}' \leftarrow$ closest vectors in $B\mathbb{Z}^n$ to c.
    if $\|v' - c\| = r$ for a $v' \in \mathcal{V}'$ **then**
      return $\mathcal{V}'$.
    end if
    $v \leftarrow$ one element of $\mathcal{V}'$.
end do

One way to speed up the convergence of this algorithm in practice is to heuristically choose an initial vector v with a sphere $S(c, r)$ of small radius.

3.5. **Checking equivalence.** We have to test equivalence and compute stabilizers under the group $\mathrm{Iso}(L)$ of Delone cells of different dimensions. Below we propose three different methods for this.

We can encode a Delone cell D by the center $c(D)$ of the empty sphere around it or by the vertex barycenter $g(D) = \frac{1}{|\operatorname{vert} D|} \sum_{v \in \operatorname{vert} D} v$ of its vertex set $\operatorname{vert} D$. Both $c(D)$ and $g(D)$ are invariant under the stabilizer of D. Any two full-dimensional Delone cells D, D' are equal if and only if $c(D) = c(D')$. However, it is possible if $n \geq 3$ that $c(D)$ lies outside or on the boundary of D. If $c(D)$ lies on the boundary of D then a facet containing $c(D)$, which is itself a Delone cell, has the same center as D. Hence, the sphere centers can be used to distinguish full-dimensional Delone cells but they do not distinguish Delone cells. Therefore, we use the vertex barycenter.

In the first method we consider the classes of the vertex barycenters $g(D)$ and $g(D')$ in the quotient $\mathbb{R}^n/L$ and check their equivalence under the finite group $\operatorname{Iso}(L)/L \simeq \mathrm{O}(L)$. The generic methods underlying isomorphism and stabilizer computations generate the full orbit of $g(D)$ under $\operatorname{Iso}(L)/L$. This is typically memory intensive. In some cases we can use a method from computational group theory, which we now explain in an example. Suppose $g(D)$ is expressed in a basis $(b_1, \ldots, b_n)$ of L as $(\frac{\alpha_1}{2 \cdot 3}, \ldots, \frac{\alpha_n}{2 \cdot 3})$ with $0 \leq \alpha_i \leq 5$ and we want to compute its stabilizer under the group $\operatorname{Iso}(L)/L$. The vector $2g(D)$ reduced modulo L is expressed as $(\frac{\tilde{\alpha}_1}{3}, \ldots, \frac{\tilde{\alpha}_n}{3})$ with $0 \leq \tilde{\alpha}_i \leq 2$. We first compute the stabilizer H of the vector $2g(D)$ under the action of $\operatorname{Iso}(L)/L$. The stabilizer of $g(D)$ under $\operatorname{Iso}(L)/L$ is equal to the stabilizer of $g(D)$ under H. This method generalizes to more than two prime factors and it is more memory efficient because the generated orbits are smaller.

The second method uses finite metric spaces of the vertex set of full-dimensional Delone cells obtained from the metric $\|v - v'\|^2$ [DL97, Chapter 14]. A finite metric space defines an edge-weighted graph. Testing if two edge weighted graphs are isomorphic can be reduced to testing if two vertex-weighted graphs are isomorphic (see [McK06, Page 25]). In practice, the program `nauty` [McK06] can solve the isomorphism problem if the number of vertices of D and D' is not too large. If the metric spaces are not isomorphic, then D and D' are not equivalent under $\operatorname{Iso}(L)$. If they are isomorphic then every graph isomorphism corresponds to a linear isometry between D and D' [BDS07, DSV07]. For each of those isomorphisms, we check if it belongs to $\operatorname{Iso}(L)$. This method is useful when the isometry group of D is small.

For the third method, we use laminations over the n-dimensional lattice L. Let D be a Delone cell of L with vertex barycenter g. One defines an $(n + 1)$-dimensional lattice $L(g)$ by taking L and adding another layer $L + v$ such that $v - g$ is orthogonal to L and $\|v - g\| = 1$. If u is an element of $\mathrm{O}(L(g))$ preserving every layer of the lamination, then it maps the vector v to some vector $w = v + h$ with $h \in L$. The function $x \mapsto u(x) + h$ preserves the Delone cell and every element preserving the layers is obtained in this way. In practice, we can use the program `AUTO` (see [PP97]) of the package `CARAT` (see [OPS98]) for computing this automorphism group. The isomorphism problem is treated similarly using the program `ISOM`.

3.6. Working with Gram matrices and periodic structures.
In many cases, it is more convenient to work with the Gram matrix $B^T B$ instead of the lattice basis B

(see [CS99, Chapter 2.2]). For instance, when B is irrational but $B^T B$ is rational. Note that our algorithms can be reformulated in terms of Gram matrices. Note also that all our algorithms can be modified to deal with periodic point sets, that is for finite unions of lattice translates. Our implementation is available from [Dut07].

3.7. **Comparison.** In [VB96] Viterbo and Biglier describe another algorithm for computing the Voronoi cell of a lattice, called the diamond cutting algorithm. As in our approach they start with a parallelepiped P defined by the basis vectors. Then they determine all lattice vectors which lie in a sphere containing $2P$. This set contains all facet defining lattice vectors of $V(L)$. Successively they add cutting planes obtained from these vectors and update the complete face lattice of the tentative Voronoi cell. They terminate when its volume coincides with $\det L$. Their implementation uses floating point arithmetic.

In comparison, our approach has the following advantages: We use the presence of symmetry in an efficient way. We do not need to compute a huge initial list of potential facet defining lattice vectors. Our algorithm does not need to compute the face lattice, not even for computing the quantizer constant as explained below. Our implementation uses rational arithmetic only.

4. COMPUTING QUANTIZER CONSTANTS

Recall from the introduction that the quantizer constant of a lattice L is the integral

$$G(L) = (\det L)^{-(1+2/n)} \int_{V(L)} \|x\|^2 dx.$$

A standard method for computing the integral $G(L)$ is to decompose $V(L)$ into simplices. Suppose that S is a simplex with vertices $v_1, \ldots, v_{n+1}$ in $\mathbb{R}^n$. Then (see [CS99, Chapter 21, Theorem 2]) the following holds:

$$\int_S \|x\|^2 dx = \frac{\text{vol } S}{(n+1)(n+2)} \left(\left\| \sum_{i=1}^{n+1} v_i \right\|^2 + \sum_{i=1}^{n+1} \|v_i\|^2 \right).$$

Thus $G(L)$ can be obtained by summing the integrals of all simplices in a decomposition of $V(L)$. Several practical methods for decomposing a polytope into simplices are discussed in [BEF98]. In our implementation, we use the triangulation obtained by the program `lrs`. However, this method as well as the other methods explained in [BEF98] are sometimes impractical and they do not use symmetries.

In order to get a group invariant subdivision, we can use the barycentric subdivision of P. That is, given any flag $F_0 \subset F_1 \subset \ldots \subset F_n = P$ of faces of P, we associate the simplex with vertex set $g_0, g_1, \ldots, g_n$ where g_i is the vertex barycenter of F_i. Note that in general there is a difference between the barycenter $\frac{1}{\text{vol } P} \int_P x dx$ of a polytope P and its vertex barycenter $\frac{1}{|\text{vert } P|} \sum_{v \in \text{vert } P} v$. The group acts on the barycentric subdivision and the stabilizer of each simplex is trivial. In practice, the number of orbits of flags can be too large.

We propose a hybrid approach, which combines the benefits of both methods. Let $\mathcal{F}$ be the facet set of an n-dimensional polytope P. We can assume without loss of generality that P has the origin as its vertex barycenter. We then have

$$(4) \qquad \int_P \|x\|^2 dx = \sum_{F \in \mathcal{F}} \int_{\mathrm{conv}(F,0)} \|x\|^2 dx.$$

To compute this sum, it is sufficient to compute the integrals only for orbit representatives of facets. Let F be a facet of P and p_F be a point in the affine space spanned by F. Then we have

$$\int_{\mathrm{conv}(F,0)} \|x\|^2 dx = \frac{1}{n+2} \left(\int_F \|y - p_F\|^2 dy \right.$$
$$\left. + 2 \int_F (y - p_F) \cdot p_F \, dy + \mathrm{vol}\, F \|p_F\|^2 \right).$$

If p_F is the orthogonal projection of the origin 0 onto F then the second summand vanishes. This point may not be invariant under the automorphism group of the facet F, but the vertex barycenter is. If we use the vertex barycenter, we also have to compute the barycenter of the polytope F as well as the volume and the square integral. In order to use symmetries coming from non-orthogonal linear transformations of P, we use the matrix valued integral

$$I_{0,1,2}(P) = \int_P \begin{pmatrix} 1 \\ x \end{pmatrix} (1, x^t) \, dx.$$

This integral splits according to

$$I_{0,1,2}(P) = \begin{pmatrix} I_0(P) & I_1(P)^t \\ I_1(P) & I_2(P) \end{pmatrix},$$

where

$$I_0(P) = \int_P dx = \mathrm{vol}\, P, \quad I_1(P) = \int_P x \, dx, \quad I_2(P) = \int_P x x^t \, dx.$$

Let G be a group of automorphisms of P. Let $O_1, \ldots, O_r$ be the G-orbits of facets of P, with representatives $F_1, \ldots, F_r$. Then the integral $I_{0,1,2}(P)$ simplifies to

$$I_{0,1,2}(P) = \sum_{i=1}^r |O_i| \left(\frac{1}{|G|} \sum_{g \in G} g I_{0,1,2}(\mathrm{conv}(F_i, 0)) g^t \right).$$

Note here that the action of the linear transformations in G is extended to homogeneous coordinates.

Assume that $I_{0,1,2}(\mathrm{conv}(F_i, 0))$ is already computed. To compute the sum in the parenthesis, we first incrementally compute a basis of the affine hull of the orbit $\{g I_{0,1,2}(\mathrm{conv}(F_i, 0)) g^t : g \in G\}$. The only G-invariant element of the affine hull is the sum we want to compute.

We now want to compute $I_{0,1,2}(\mathrm{conv}(F, 0))$ in terms of lower dimensional integrals. The integral depends on the chosen basis. If f is an affine transformation of

$\mathbb{R}^n$, then we have the following transformation formula

$$f I_{0,1,2}(fP) f^t |\det f| = I_{0,1,2}(P),$$

for any an n-dimensional polytope P in $\mathbb{R}^n$. This allows to compute $I_{0,1,2}(P)$ for another basis. So, we can choose a coordinate system such that

$$F = \left\{ \begin{pmatrix} 1 \\ x \end{pmatrix} : x \in F' \right\} \subset \mathbb{R}^n,$$

where $F' \subset \mathbb{R}^{n-1}$ is an $(n-1)$-dimensional polytope. We then have the following formulas:

$$I_0(\text{conv}(F, 0)) = \tfrac{1}{n} I_0(F'), \quad I_1(\text{conv}(F, 0)) = \tfrac{1}{n+1} \begin{pmatrix} I_0(F') \\ I_1(F') \end{pmatrix},$$
$$I_2(\text{conv}(F, 0)) = \tfrac{1}{n+2} I_{0,1,2}(F').$$

For computing $I_{0,1,2}(F')$, we have two options: Either we use the first method of this section, which involves computing a triangulation or we apply the above method recursively. The decision is made heuristically, depending on the size of the automorphism group of F and its number of vertices. In order to reduce the size of the computation, one can store intermediate results.

Those methods are general and apply to any polytope and any polynomial function, which we want to integrate over P. Note that a similar method of using the standard formula (4) has been used for computing the volume in [BEF98] under the name of *Lasserre's method* ([Las98]), albeit in a non-group setting.

5. RESULTS

In this section, we collect results from our implementation of the algorithms explained in Sections 3 and 4. We obtain previously unknown exact covering densities and quantizing constants of several prominent lattices and their duals. Recall that the *dual* L^* of a lattice $L \subset \mathbb{R}^n$ is defined by

$$L^* = \{x \in \mathbb{R}^n : y \cdot x \in \mathbb{Z} \quad \text{for all } y \in L\}.$$

The *covering density* of an n-dimensional lattice L is

$$\frac{\mu(L)^n}{\det L} \operatorname{vol} B_n,$$

where B_n is the unit ball in $\mathbb{R}^n$. Other computations of Voronoi cells of lattices can be found in [CS91], [EMS03, Chapter 5] and [MP95].

5.1. Coxeter lattices. The root lattice A_n is defined by

$$\mathsf{A}_n = \{x \in \mathbb{Z}^{n+1} : \sum_{i=1}^{n+1} x_i = 0\}.$$

If r divides $n + 1$, the *Coxeter lattice* A_n^r (see [Cox51]) is defined by translates of A_n:

$$\mathsf{A}_n^r = \mathsf{A}_n \cup (v_n^r + \mathsf{A}_n) \cup \ldots \cup ((r-1)v_n^r + \mathsf{A}_n),$$

where $v_n^r = \frac{1}{n+1} \sum_{i=2}^{n+1} (e_i - e_1)$. The dual lattice of A_n^r is $\mathsf{A}_n^{n+1/r}$.

The Delone decomposition of the lattice A_n^r has been studied in [Anz02, Anz06, Bar94] up to dimension $n = 15$, hoping to obtain lattices with low covering density. One pleasant fact is that the symmetry group of A_n^r contains the group $\mathrm{Sym}(n + 1) \times \mathbb{Z}_2$. Latter can be represented as a permutation group acting on $n + 3$ points, which drastically simplifies isomorphism computations.

In Table 1 we list the obtained results. Note that the lattices A_{17}^6, A_{19}^{10}, A_{20}^7 and A_{21}^{11} turn out to give new record sphere coverings. Up to dimension 8 all those lattices are well known and their Voronoi cells can be obtained by standard computer algebra software. Our list is complete up to dimension 21. For the missing cases we could not finish the computation.

5.2. **Laminated lattices.** The laminated lattices, which are defined in [CS99, Chapter 6], give the best known lattice sphere packings in many dimensions. The Delone subdivision is known up to dimension 8 and in dimension 24 for all laminated lattices and their duals [CS99, Chapters 21, 23, 25]. In dimension 16, the covering density of Λ_{16} is known [CS99, Chapter 6].

In Table 2 we list the obtained results, which are complete up to dimension 17.

5.3. **Shorter Leech lattice.** The 4600 shortest vectors of Λ_{23}^* of define a sublattice of index 2, called the *shorter Leech lattice* O_{23} ([CS99, Page 179, 420, 441]).

The Delone decomposition (see Table 3) is remarkable in many respects: There are only 5 orbits and the first one has the full symmetry group of the lattice. It turns out that $\Lambda_{23}^* = O_{23} \cup (v + O_{23})$ where v is the center of a centrally symmetric Delone cell lying in the first orbit. The covering density of O_{23} is 15218.062669.

5.4. **Cut lattices.** The cut polytope CUT_n is a famous polytope appearing in combinatorial optimization (see [DL97]). It has 2^{n-1} vertices and is of dimension $\frac{n(n-1)}{2}$. The lattice generated by its vertices is called *cut lattice* and is denoted by $L(\mathrm{CUT}_n)$ (see [DG95]). The polytope CUT_n is one of its full-dimensional Delone cells. We list our results in Table 4.

5.5. **Quantizer constants.** In Table 5 we collect some new exact quantizer constants.

According to [AE98], the lattice D_{10}^+ is conjectured to be the optimal lattice quantizer. Conway and Sloane approximated $G(K_{12})$ ([CS99, Table 2.3]) using Monte-Carlo integration; our exact computation fits into their bounds.

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lattice	# orbits	covering density	lattice	# orbits	covering density
A_9^2	6	18.543333	A_9^5	5	4.340184
A_{11}^2	6	94.090996	A_{11}^3	11	27.089662
A_{11}^4	16	5.598337	A_{11}^6	4	7.618558
A_{13}^2	10	134.623484	A_{13}^7	10	7.864060
A_{14}^3	17	32.313517	A_{14}^5	31	9.006610
A_{15}^2	10	722.452642	A_{15}^4	19	25.363859
A_{15}^8	10	11.601626			
A_{17}^2	15	1068.513081	A_{17}^3	26	240.511580
A_{17}^6	73	12.357468	A_{17}^9	24	17.231927
A_{19}^2	15	5921.056764	A_{19}^4	58	40.445924
A_{19}^5	80	25.609662	A_{19}^{10}	80	21.229200
A_{20}^3	40	307.209487	A_{20}^7	187	20.366828
A_{21}^2	21	8937.567486	A_{21}^{11}	64	27.773140
A_{23}^3	55	2405.032746	A_{23}^4	85	205.561225
A_{23}^6	187	79.575330	A_{23}^{12}	100	43.231587
A_{24}^5	144	115.011591			
A_{25}^{13}	210	54.472182			
A_{26}^3	75	3184.1387034	A_{26}^9	1231	50.937168
A_{27}^4	156	350.137031	A_{27}^{14}	338	76.909712
A_{29}^3	102	25664.644103	A_{29}^5	347	202.040331
A_{29}^6	711	154.329831	A_{29}^{10}	3581	84.324725
A_{29}^{15}	678	114.084219			
A_{31}^{16}	1225	33.934941			

TABLE 1. Number of orbits of full-dimensional Delone cells and covering density for some Coxeter lattices

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lattice	# orbits	covering density	lattice	# orbits	covering density
Λ_9	5	9.003527	Λ_9^*	9	9.003527
Λ_{10}	7	12.408839	Λ_{10}^*	21	9.306629
Λ_{11}^{max}	11	24.781167	Λ_{11}^{max*}	18	19.243468
Λ_{11}^{min}	18	24.781167	Λ_{11}^{min*}	153	8.170432
Λ_{12}^{max}	5	30.418954	Λ_{12}^{max*}	8	42.728408
Λ_{12}^{mid}	23	30.418954	Λ_{12}^{mid*}	52	19.176309
Λ_{12}^{min}	13	30.418954	Λ_{12}^{min*}	78	12.292973
Λ_{13}^{max}	18	60.455139	Λ_{13}^{max*}	57	43.214494
Λ_{13}^{mid}	46	35.931846	Λ_{13}^{mid*}	125	19.155991
Λ_{13}^{min}	129	60.455139	Λ_{13}^{min*}	5683	13.724864
Λ_{14}	65	98.875610	Λ_{14}^*	1392	34.721750
Λ_{15}	27	202.910873	Λ_{15}^*	108	25.642067
Λ_{16}	4	96.500266	Λ_{16}^*	4	96.500266
Λ_{17}	28	197.719499	Λ_{17}^*	720	100.173101
Λ_{18}	239	301.192334			

TABLE 2. Number of orbits of full-dimensional Delone cells and covering density for some laminated lattices and their duals

number of vertices	size of stabilizer group
94208	84610842624000
32	1344
24	10200960
24	1320
24	1320

TABLE 3. Orbits of full-dimensional Delone cells of O_{23}

lattice	dimension	# orbits	covering density
$L(\text{CUT}_3)$	3	2	2.09439
$L(\text{CUT}_4)$	6	4	5.16771
$L(\text{CUT}_5)$	10	12	40.80262
$L(\text{CUT}_6)$	15	112	255.4255

TABLE 4. Dimensions, number of orbits of full-dimensional Delone cells and covering density of some cut lattices

lattice	quantizer constant
Λ_9	$\frac{151301}{2099520} \approx 0.07206$
Λ_9^*	$\frac{1371514291}{19110297600} \approx 0.07176$
A_9^2	$\frac{2120743}{\sqrt[9]{5.2^8 13271040}} \approx 0.072166$
A_9^5	$\frac{8651427563}{\sqrt[9]{2.5^8 26578125000}} \approx 0.072079$
D_{10}^+	$\frac{4568341}{64512000} \approx 0.07081$
A_{11}^2	$\frac{452059}{\sqrt[11]{35702400}} \approx 0.07174$
A_{11}^3	$\frac{287544281699}{\sqrt[11]{4.3^{10} 1325839006800}} \approx 0.070426$
A_{11}^4	$\frac{6387657954959}{\sqrt[11]{3.2^9 46506442752000}} \approx 0.070494$
D_{12}^+	$\frac{29183629}{412776000} \approx 0.070700$
K_{12}	$\frac{797361941}{\sqrt{36567561000}} \approx 0.070095$

TABLE 5. Quantizer constants of some lattices

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